

Generic Anonymity Wrapper for Messaging Protocols

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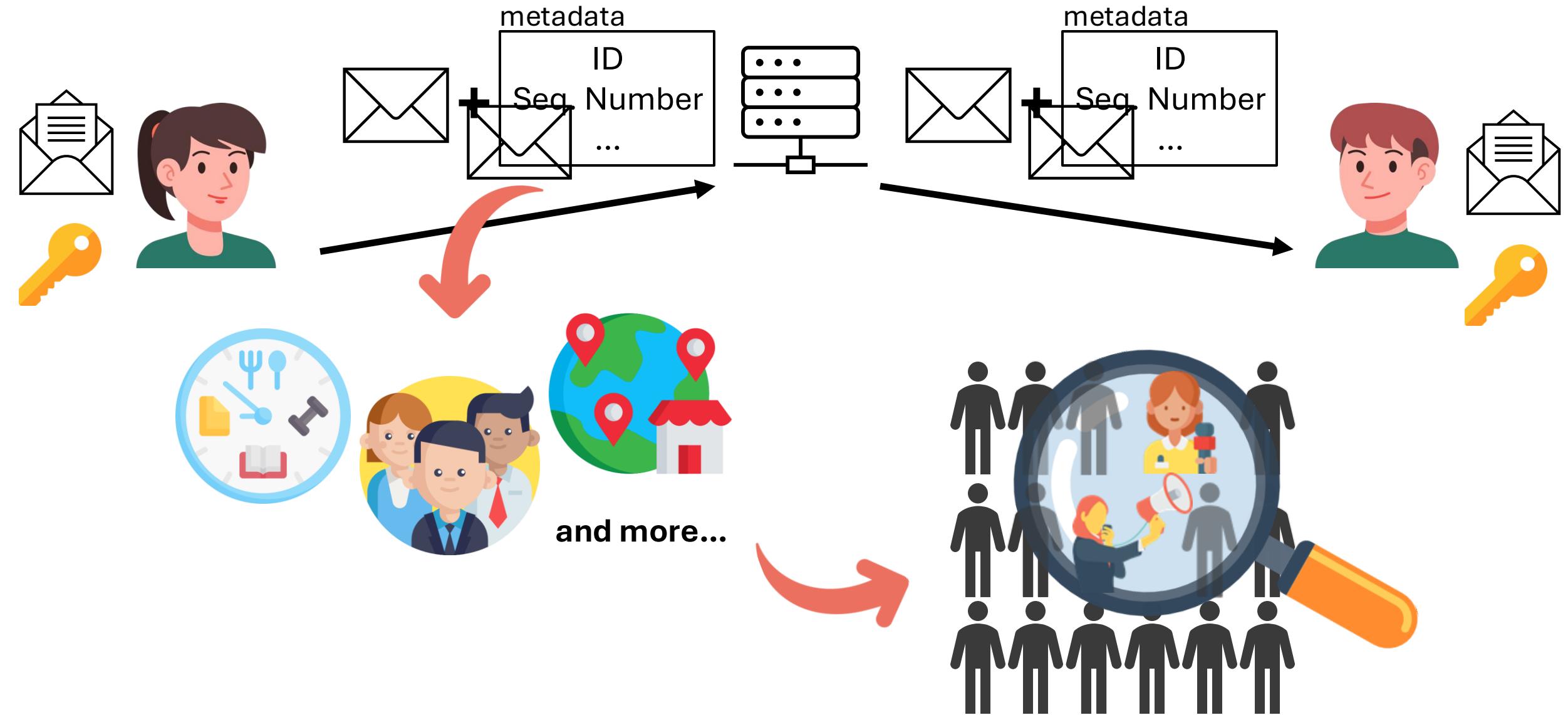
Alexander Bienstock

J.P. Morgan AI Research and J.P.
Morgan AlgoCRYPT CoE

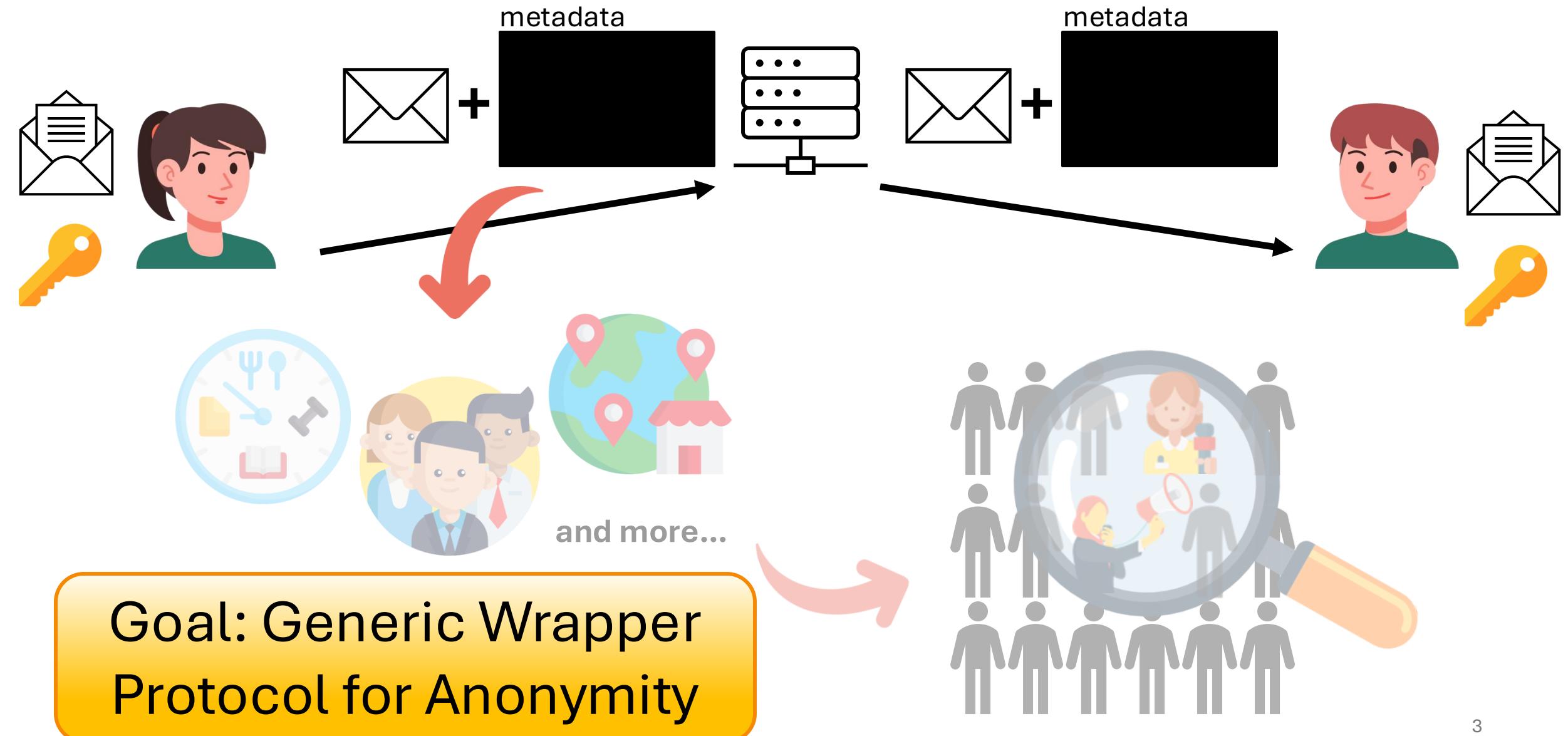
Rolfe Schmidt
Signal Messenger

Yevgeniy Dodis
New York University

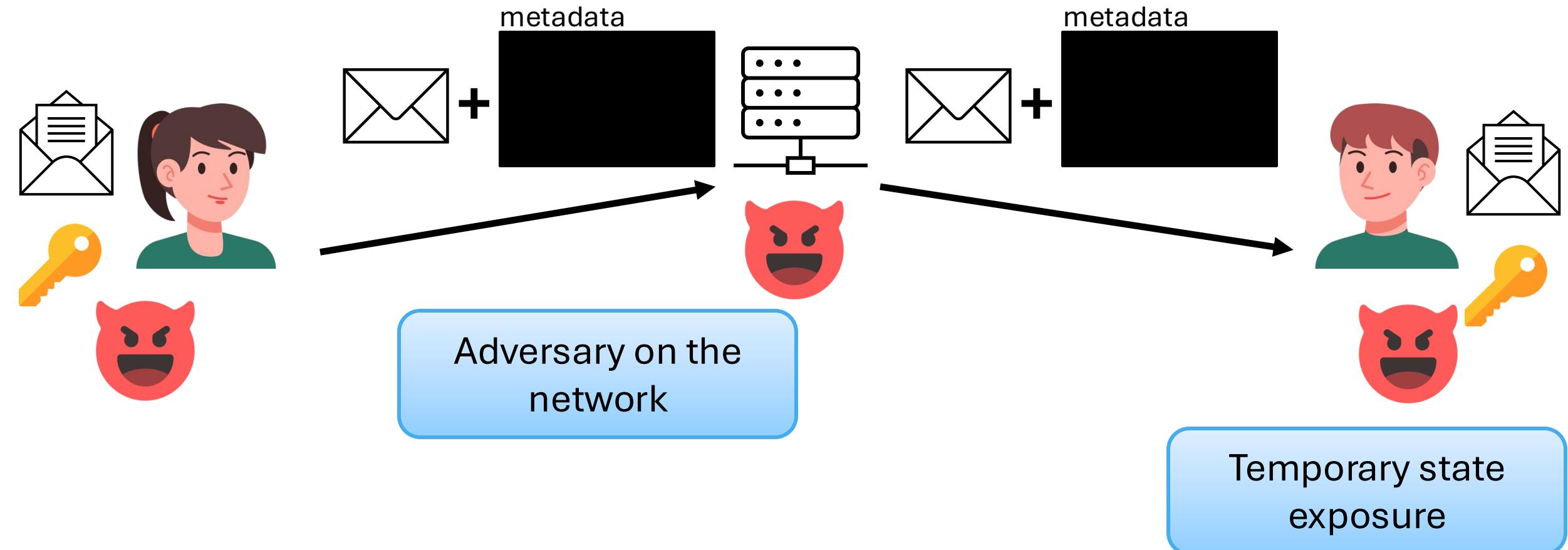
Metadata in Messaging Protocols



Metadata in Messaging Protocols

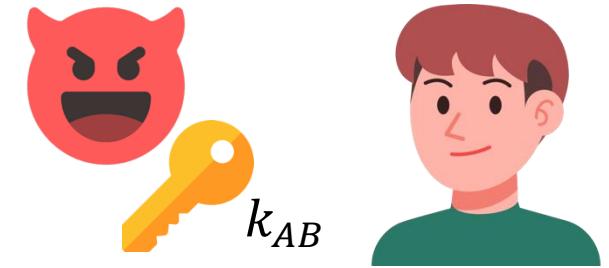
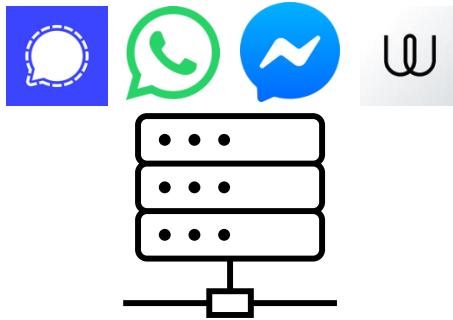


Threat Model

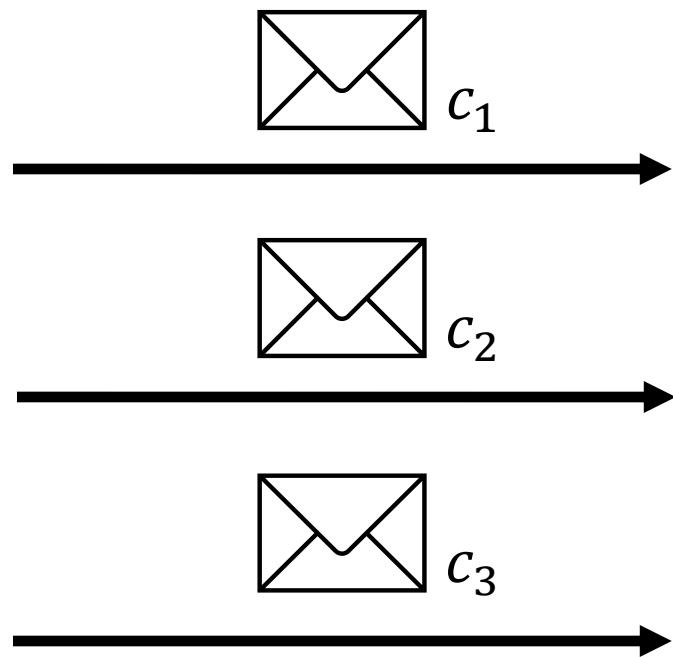


Goal: Generic Wrapper
Protocol for Anonymity

Secure communication

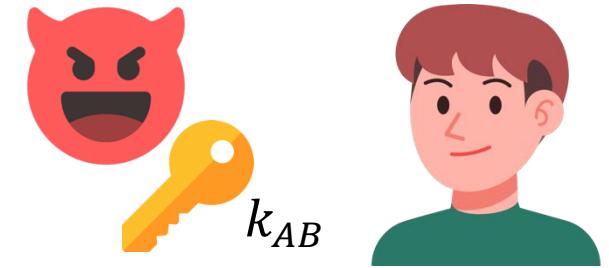
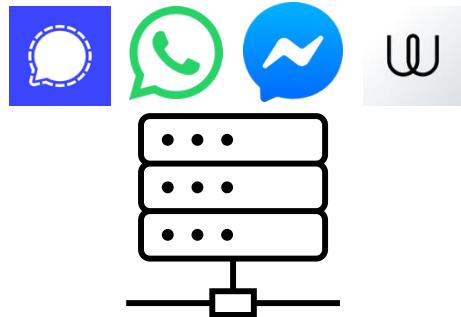


$$c_1 = aenc_{k_{AB}}(m_1)$$



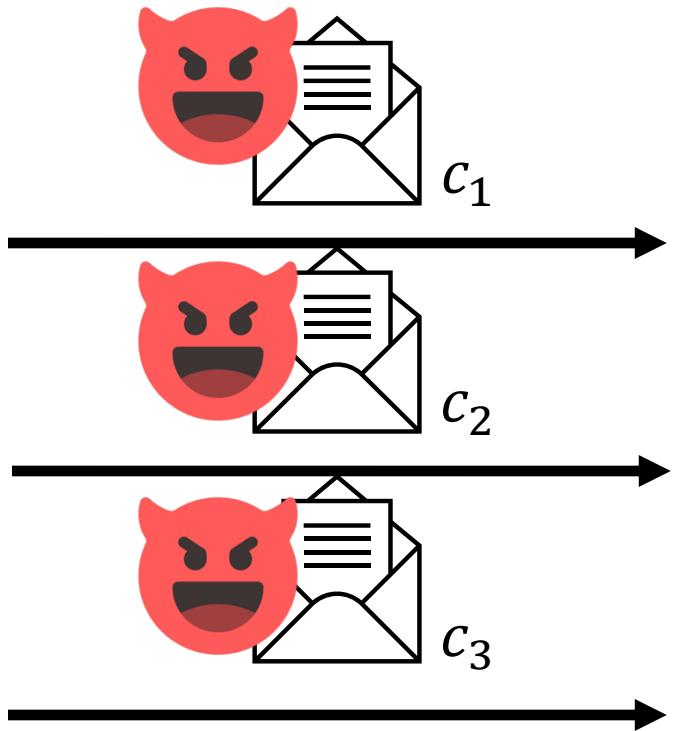
$$m_1 = adec_{k_{AB}}(c_1)$$

Secure communication

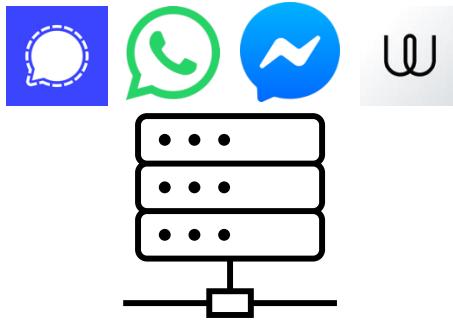


$$c_1 = aenc_{k_{AB}}(m_1)$$

$$m_1 = adec_{k_{AB}}(c_1)$$

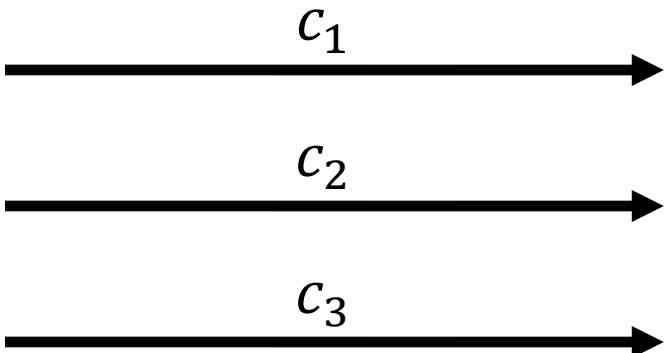


Secure communication



$$\begin{array}{c} ck_0 \\ \searrow \\ ck_1 k_1 \\ \searrow \\ ck_2 k_2 \\ \searrow \\ ck_3 k_3 \\ \dots \\ \searrow \\ ck_0^* \text{ ✨} \\ \searrow \\ ck_0^* k_0^* \end{array}$$

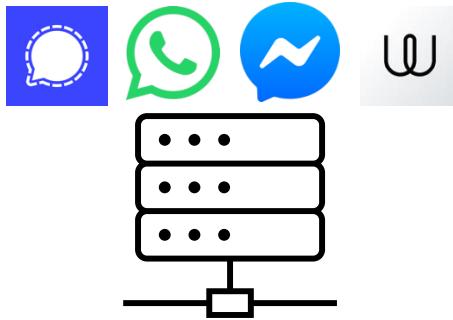
$$\begin{aligned} c_1 &= aenc_{k_1}(m_1) \\ c_2 &= aenc_{k_2}(m_2) \\ c_3 &= aenc_{k_3}(m_3) \end{aligned}$$



$$\begin{aligned} m_1 &= adec_{k_1}(c_n) \\ m_2 &= adec_{k_2}(c_2) \\ m_3 &= adec_{k_3}(c_3) \end{aligned}$$

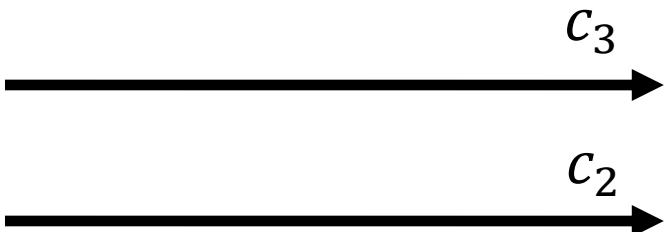
$$\begin{array}{c} ck_0 \\ \searrow \\ ck_1 k_1 \\ \searrow \\ ck_2 k_2 \\ \searrow \\ ck_3 k_3 \\ \dots \\ \searrow \\ ck_0^* \text{ ✨} \\ \searrow \\ ck_0^* k_0^* \end{array}$$

Secure communication



$$\begin{array}{c} ck_0 \\ \searrow \\ ck_1 k_1 \\ \searrow \\ ck_2 k_2 \\ \searrow \\ ck_3 k_3 \\ \dots \\ \searrow \\ ck_0^* \text{ ✨} \\ \searrow \\ ck_0^* k_0^* \end{array}$$

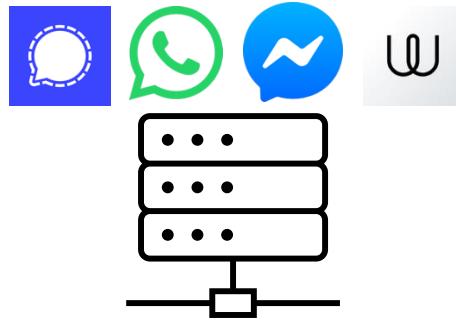
$$\begin{aligned} c_1 &= aenc_{k_1}(m_1) \\ c_2 &= aenc_{k_2}(m_2) \\ c_3 &= aenc_{k_3}(m_3) \end{aligned}$$



$$\begin{aligned} m_1 &= adec_{k_1}(c_n) \\ m_2 &= adec_{k_2}(c_2) \\ m_3 &= adec_{k_3}(c_3) \end{aligned}$$

$$\begin{array}{c} ck_0 \\ \searrow \\ ck_1 k_1 \\ \searrow \\ ck_2 k_2 \\ \searrow \\ ck_3 k_3 \\ \dots \\ \searrow \\ ck_0^* \text{ ✨} \\ \searrow \\ ck_0^* k_0^* \end{array}$$

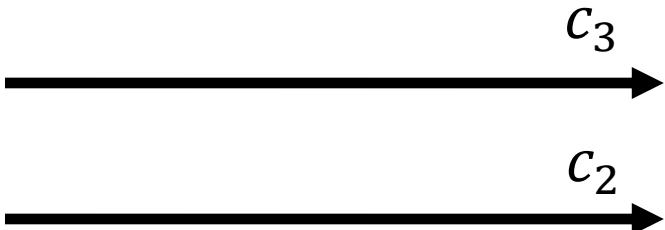
Secure communication



$$\begin{array}{l} ck_0 \\ \downarrow \\ ck_1 \ k_1 \\ \downarrow \\ ck_2 \ k_2 \\ \downarrow \\ ck_3 \ k_3 \end{array}$$

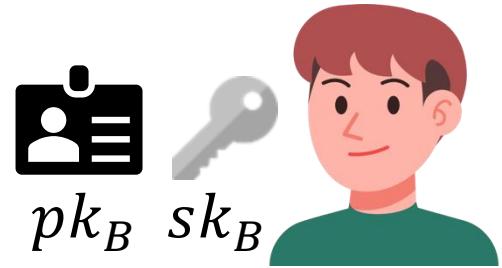
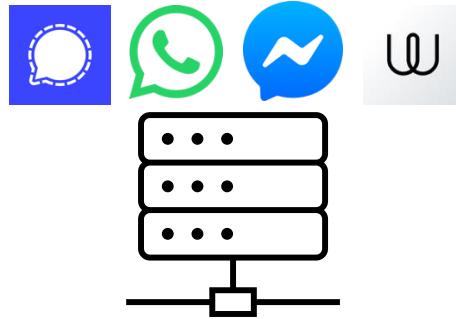
$$\begin{aligned} c_1 &= aenc_{k_1}(m_1) \\ c_2 &= aenc_{k_2}(m_2) \\ c_3 &= aenc_{k_3}(m_3) \end{aligned}$$

keep skipped keys
for immediate
decryption



$$\begin{array}{l} ck_0 \\ \downarrow \\ ck_1 \ k_1 \\ \downarrow \\ ck_2 \ k_2 \\ \downarrow \\ ck_3 \ k_3 \end{array}$$

Double Ratchet: Metadata

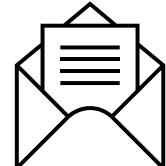


ck_0
 $ck_1 k_1$
 $ck_2 k_2$
 $ck_3 k_3$
...

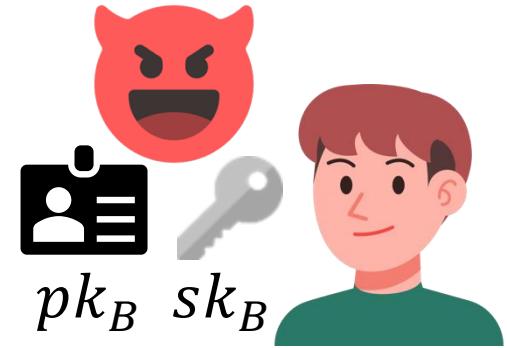
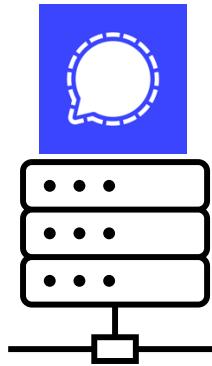
$$c_3 = aenc_{k_3}(m_3)$$

$$c_3 + \boxed{\text{index}=3, pk_A}$$

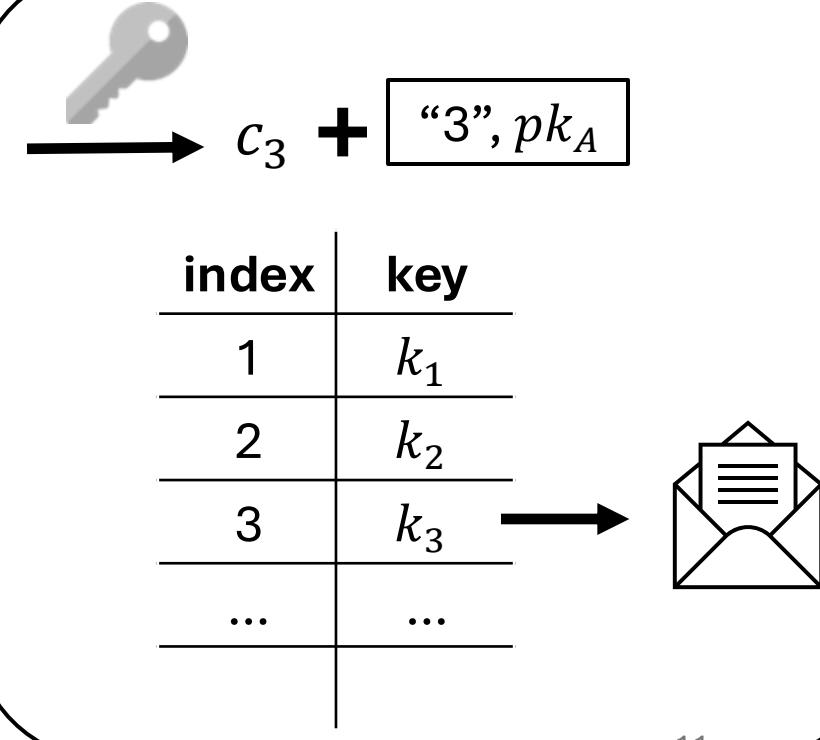
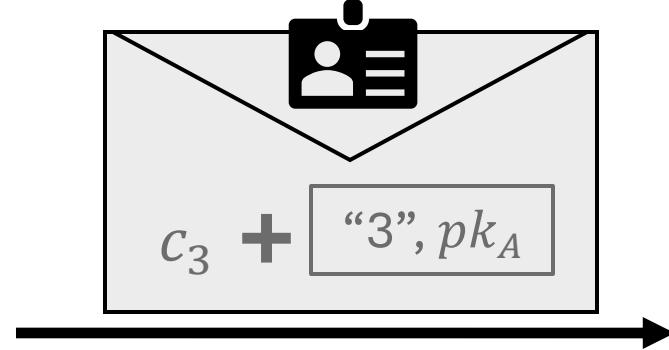
index	key
1	k_1
2	k_2
3	k_3
...	...



Hiding Metadata: Sealed Sender



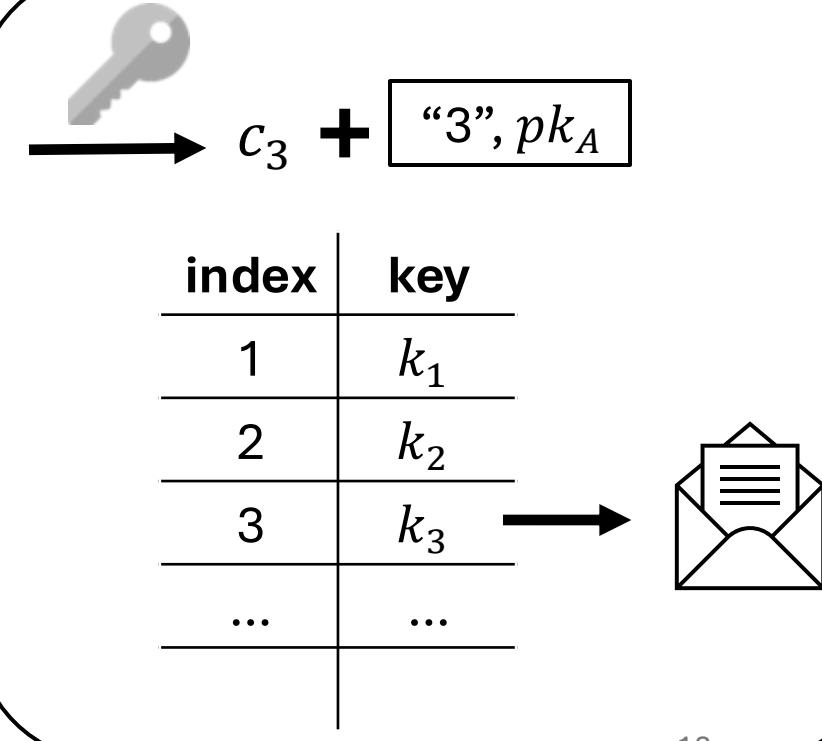
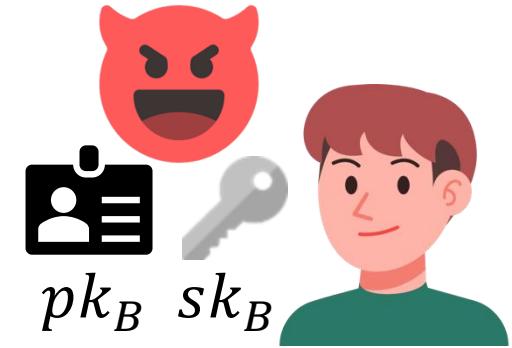
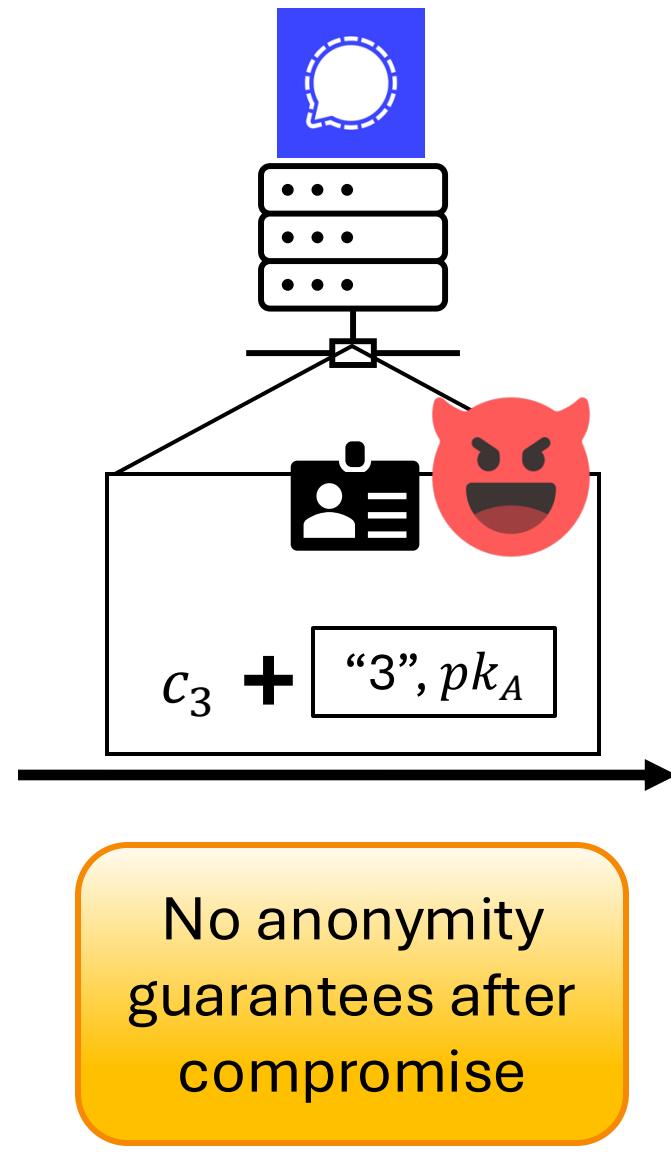
$$c_3 = aenc_{k_3}(m_3)$$



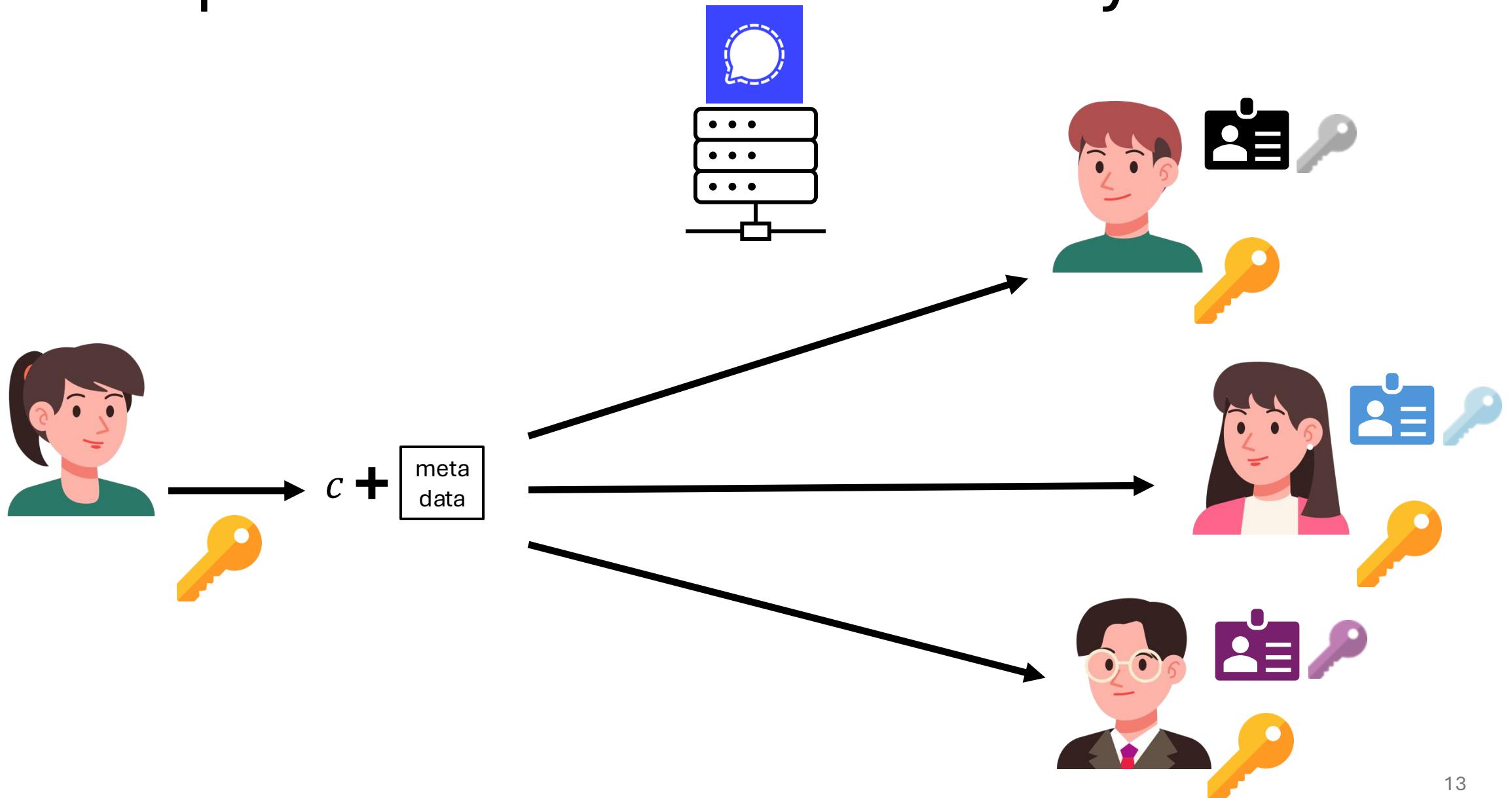
Hiding Metadata: Sealed Sender



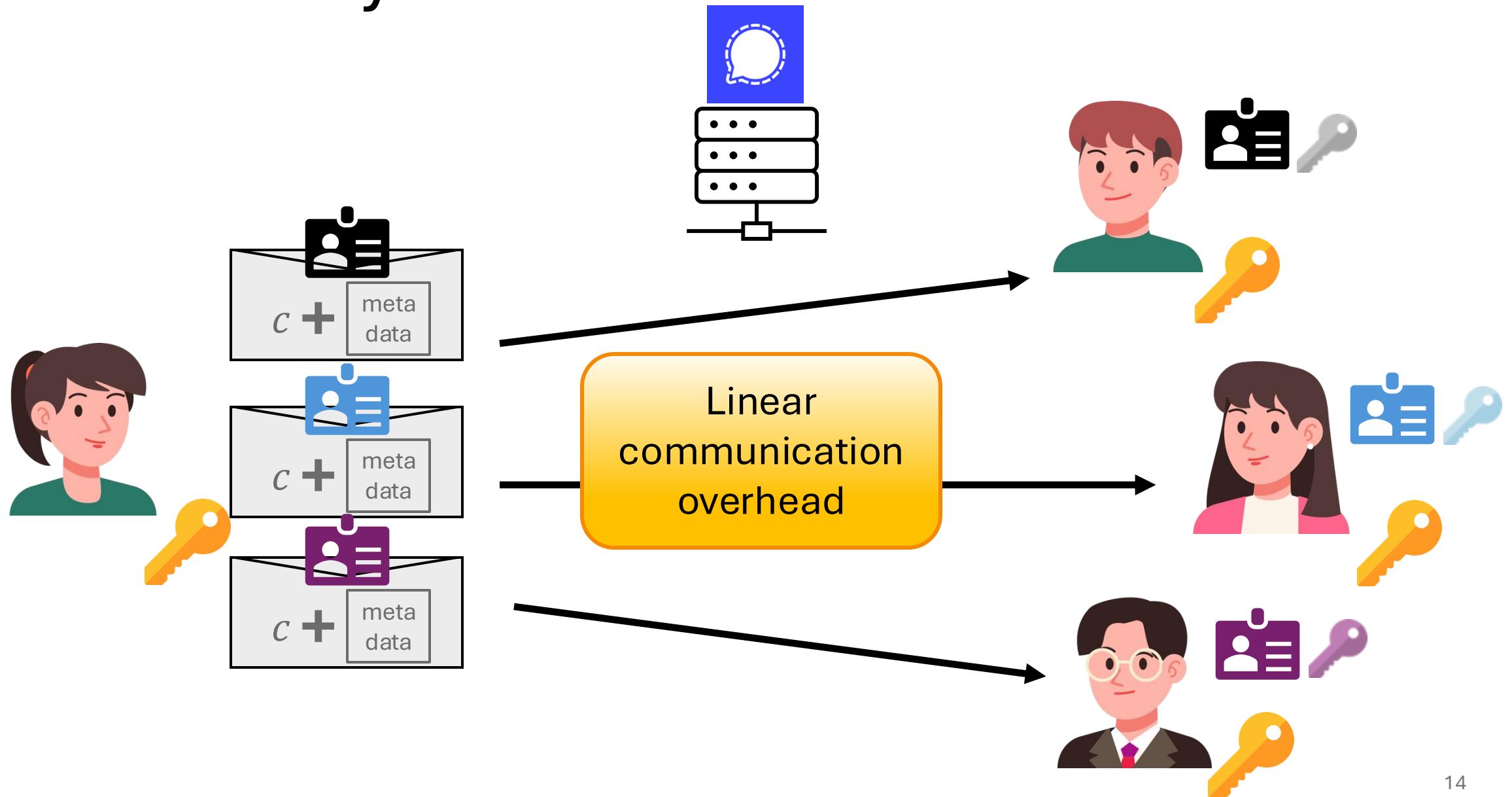
$$c_3 = aenc_{k_3}(m_3)$$



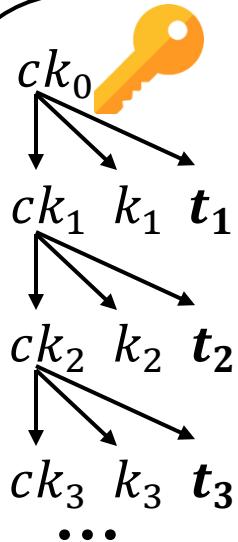
Group Communication: Sender Keys



Sender Keys + Sealed Sender



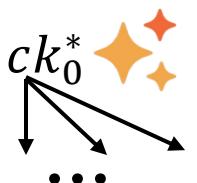
Anonymity Wrapper



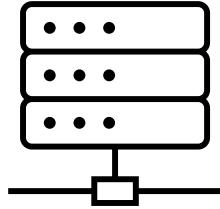
Instead of Sender Keys
+ Sealed Sender

$$c'_1 = aenc_{k_1}(m_1)$$

Or as a generic
wrapper:



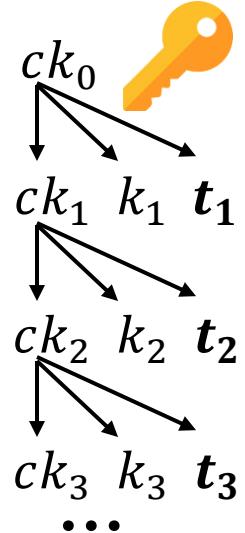
$$c'_1 = aenc_{k_1}(c_1 + \boxed{\text{meta data}})$$



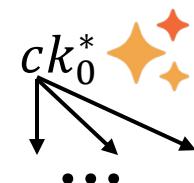
c'_1, t_1



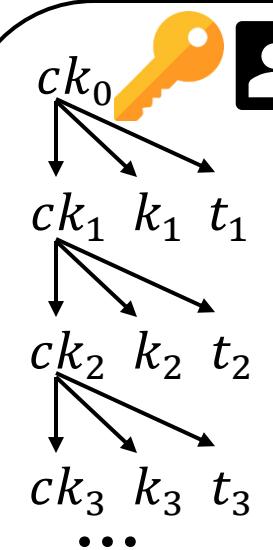
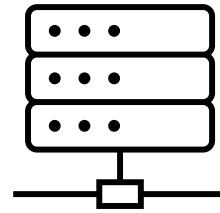
tag	key
t_1	k_1
t_2	k_2
t_3	k_3
...	...



$$m_1 = adec_{k_1}(c_1)$$



Anonymity Wrapper: Sender Authentication



Instead of Sender Keys
+ Sealed Sender

$$c'_1 = aenc_{k_1}(m_1)$$

Or as a generic
wrapper:

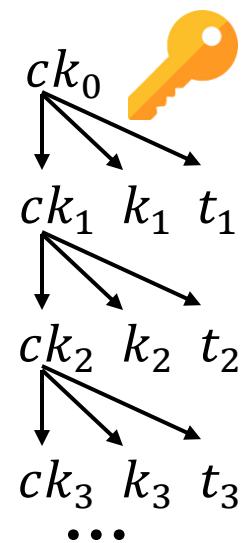
$$c'_1 = aenc_{k_1}(c_1 + \begin{matrix} \text{meta} \\ \text{data} \end{matrix})$$

c'_1, t_1, σ_1

tag	key
t_1	k_1
t_2	k_2
t_3	k_3
...	...

$\sigma_1?$

$$m_1 = adec_{k_1}(c_1)$$



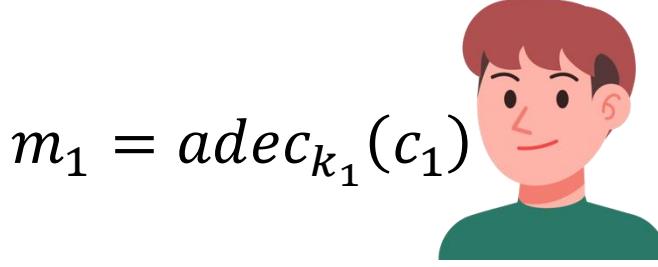
$$c'_1 = aenc_{k_1}(c_1 + \begin{matrix} \text{meta} \\ \text{data} \end{matrix})$$

Anonymity Wrapper: Performance



$$c'_1 = aenc_{k_1}(c_1 + \boxed{\text{meta data}})$$

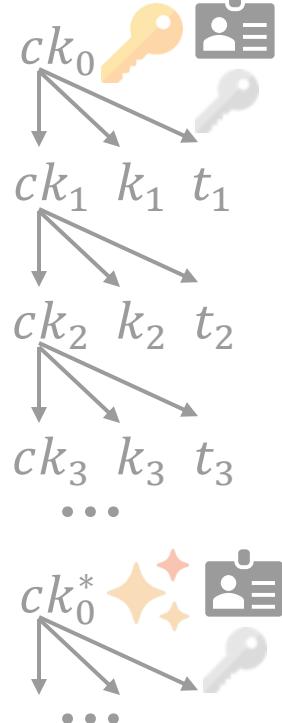
c'_1, t_1, σ_1



$$m_1 = adec_{k_1}(c_1)$$



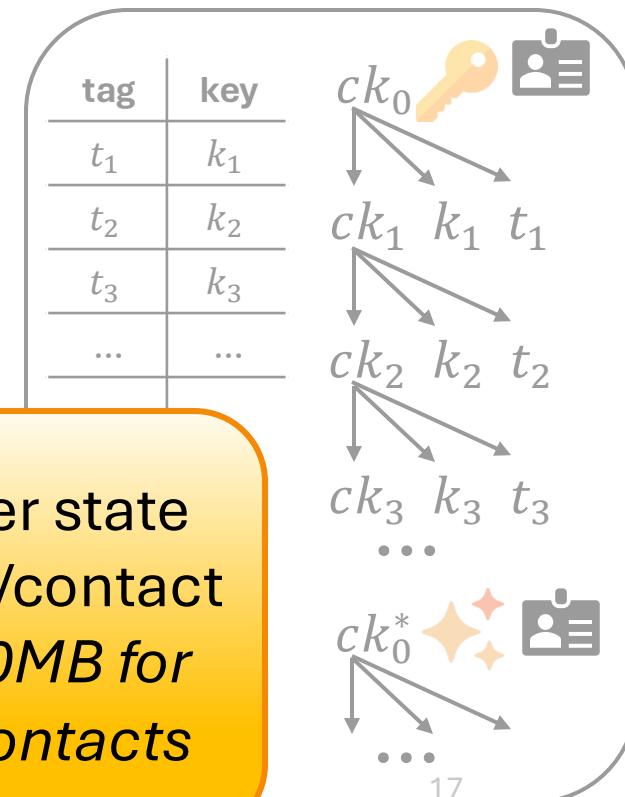
Encryption (Decryption) (μ s)					
Group Size	2	5	10	100	1000
AW					
SS+SK					



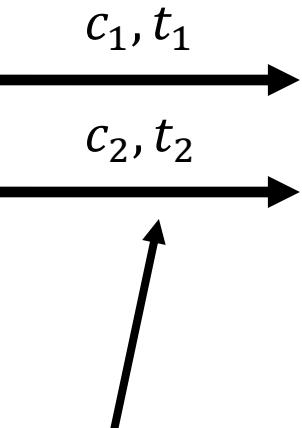
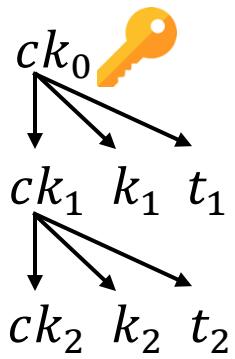
Sender state
+ 192 bytes

	Ciphertext size (bytes)
AW	
SS+SK	

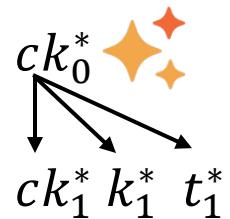
Receiver state
+ 600KB/contact
i.e., 600MB for
1000 contacts



Receiver State: Out-of-Order Delivery



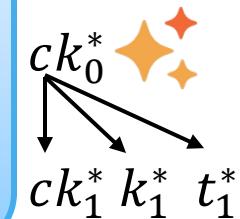
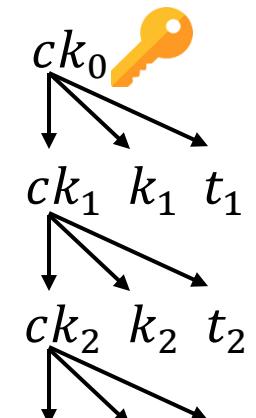
No message index to advance key chain!



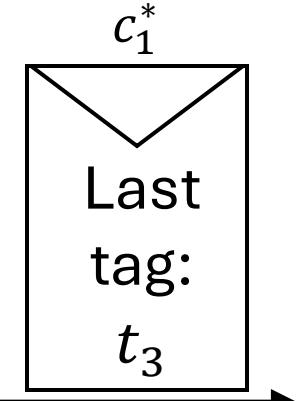
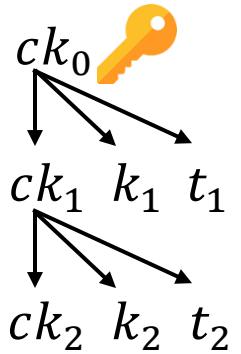
tag	key
t_1	k_1
t_2	k_2
t_3	k_3
t_4	k_4
...	...

t_2 not in receiver state
→ Continue key chain?
→ Discard...?
→ Pre-compute keys ☺

State size: Receiver stores
2000 future pre-computed
keys for each session
+ up to 2000 skipped keys

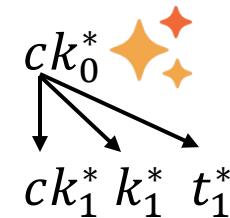
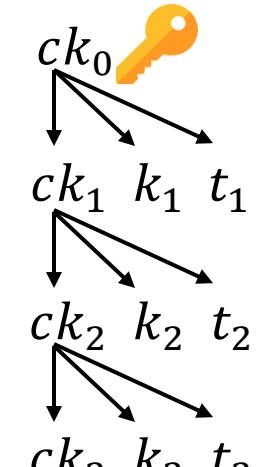
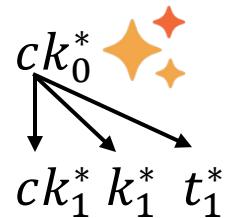


Receiver State: Pre-computed Keys

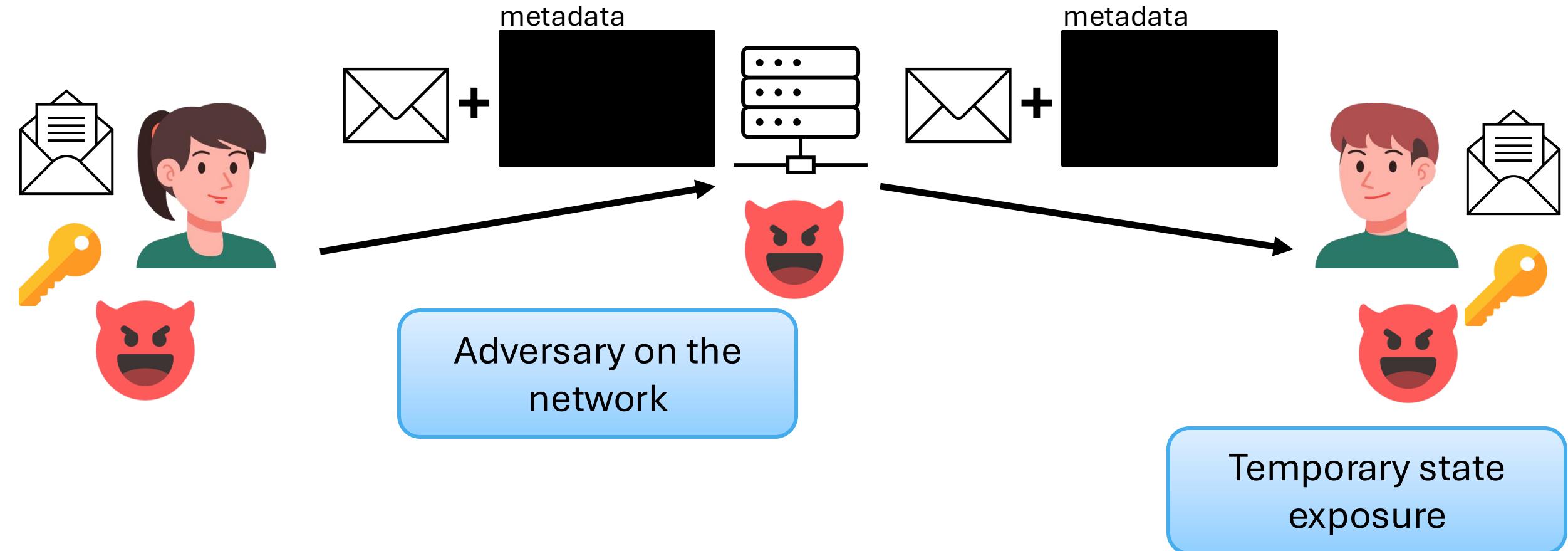


tag	key
t_2	k_2
t_3	k_3
...	...
t_{2003}	t_{2003}

← unnecessary



Threat Model



Goal: Generic Wrapper
Protocol for Anonymity

Metadata in Receiver State



- State exposure should not reveal past communication

But this state is too large, how can we improve?

tag	key
...	...

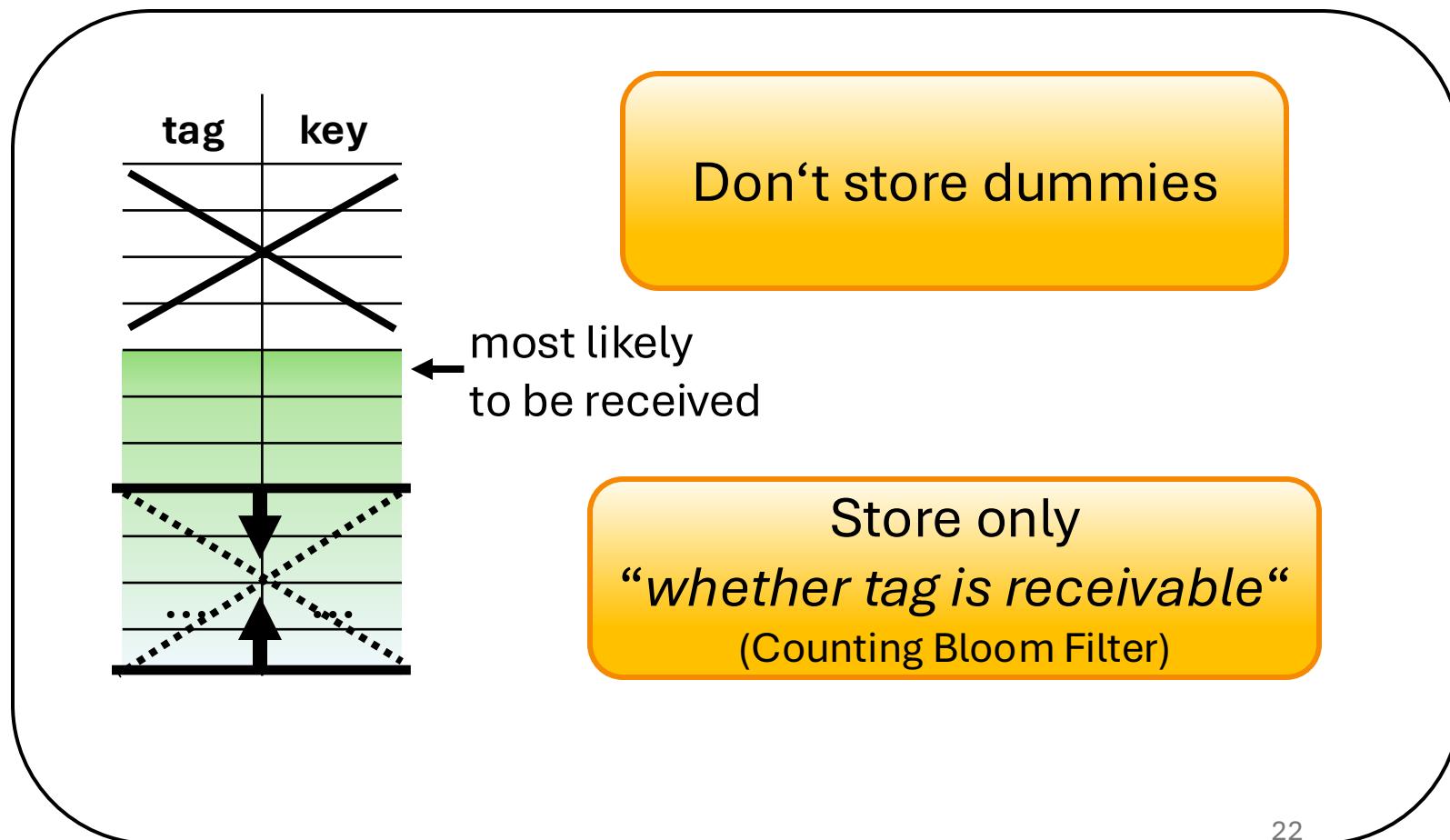
- Only store key derivation order (*no indices*)
- Constant size (may insert dummies)

Improving Size of Receiver State

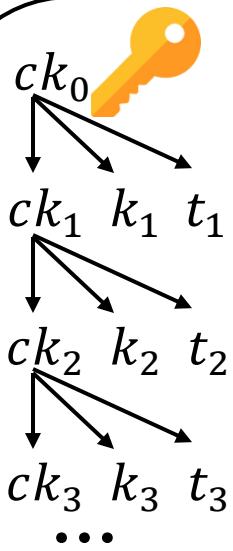


- Up to 1/3 of receiver state are dummies
- Most pre-computed keys only needed for worst case

Storage size:
38KB/contact
3.8MB/1000 contacts



Conclusion

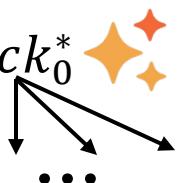


Instead of Sender Keys + Sealed Sender

$$c'_1 = aenc_{k_1}(m_1)$$

Or as a generic
wrapper:

$$c'_1 = aenc_{k_1}(c_1 + \boxed{\text{meta data}})$$



Encryption (Decryption) (μ s)					
Group Size	2	5	10	100	1000
AW	21 (32)				
SS+SK	131 (61)	283 (61)	343 (61)	1,116 (61)	6,052 (61)



	Message size (bytes)
AW	155
SS+SK	$440 + 68 * \text{group_size}$

c'_1, t_1

- Signal considers deploying this protocol

tag	key
t_1	k_1
t_2	k_2
t_3	k_3
...	...

$$m_1 = adec_{k_1}(c_1)$$

